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Analytical and Simulation Tools for Optical Camera Communications

Alexis Duque^{a,1}, Razvan Stanica^a, Herve Rivano^a, Adrien Desportes^b

^a*Univ Lyon, INSA Lyon, Inria, CITI, F-69621 Villeurbanne, France.*

^b*Rtone, Lyon, France.*

Abstract

The use of LED-to-camera communication opens the door to a wide range of use cases and applications, with diverse requirements in terms of quality of service. However, while analytical models and simulation tools exist for all the major radio communication technologies, the only way of currently evaluating the performance of a network mechanism over LED-to-camera is to implement and test it. Our work aims to fill this gap by proposing a Markov-modulated Bernoulli process to model the wireless channel in LED-to-camera communications, which is shown to closely match experimental results. Based on this model, we develop and validate *CamComSim*, the first network simulator for LED-to-camera communications.

Keywords: visible light communications, LED-to-camera, simulation tools, Markov-modulated Bernoulli process

1. Introduction

Visible-light communication (VLC) is an enabling technology that exploits illumination to provide a short-range wireless communication link. VLC systems take advantage of the license-free light spectrum and their immunity to radio frequency (RF) interference. In such systems, information is often relayed by modulating the output intensity of a light-emitting diode (LED). Any electronic device which can detect the presence or absence of visible light can be utilized as a VLC receiver. While most of the work in the field is focused towards using photo-diodes as receivers [1, 2], because of their fast response and high bandwidth, some studies demonstrated that smartphone cameras can also be used to detect high-frequency light patterns [3].

13 Indeed, nowadays smartphone cameras widely use two types of image
 14 sensors, Charge Coupled Devices (CCD) or Complementary Metal Oxide
 15 Semiconductors (CMOS). These two technologies have some similarities, but
 16 one major distinction is the way each sensor exposes its pixels to light. CCD
 17 sensors use the Global Shutter readout mode, where all pixels are exposed
 18 simultaneously and then each pixel is read sequentially. This mechanism
 19 helps in capturing a still image of a moving object. On the other hand, CMOS
 20 sensors use the Rolling Shutter readout mode [4], where each row is exposed
 21 in a row-sequential way with fixed time delay. Due to this mechanism, there
 22 is a significant time difference between the beginning of the exposure of the
 23 first and the last row, making them no longer simultaneous. When an LED
 24 is modulated at a frequency higher than the rolling shutter speed, stripes of
 25 different light intensity are captured in the image. A row of pixels appears
 26 illuminated when the LED was ON during the row exposure time. On the
 27 other hand, a row appears dark when the LED was OFF during the exposure
 28 time, as shown in Fig. 1. The intensity and width of the strip depend on
 29 the transmitter modulation frequency, allowing us to encode information in
 30 these illuminated and dark bands, similarly to the use of a bar code.

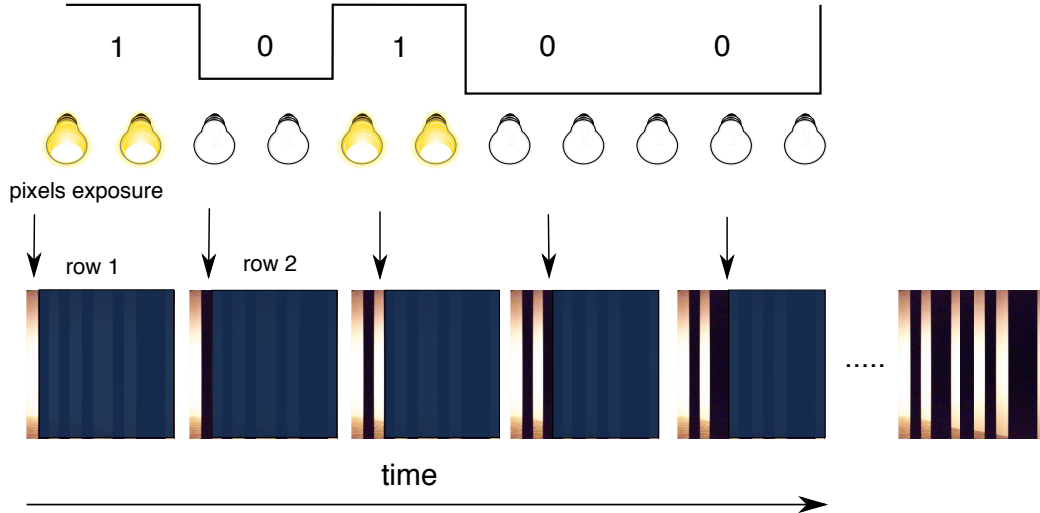


Figure 1: Stripe formation and LED state. The width of each strip corresponds to the duration of the LED state (ON or OFF).

31 This LED-to-camera communication based on the Rolling Shutter effect

opens the door to a wide range of use cases and applications, with diverse requirements in terms of quality of service [5]. To cite a few examples, both line-of-sight (LOS) [6, 7] and non-line-of-sight (NLOS) [8, 9] communications have been demonstrated in these settings, as well as ultra-reliable localization solutions [10, 11], sensing [12], or even scene protection against intrusive photographs [13]. However, while analytical models and simulation tools exist for all the major RF technologies, the only way of currently evaluating the performance of a network mechanism over LED-to-camera is to implement and test it. This results in heavy measurement and parameterisation campaigns that need to be repeated anytime a new VLC protocol or feature is imagined. Having access to standard performance evaluation tools in this type of network would certainly accelerate studies in the field, and nicely complement experimental field tests.

The work described in this paper aims to fill this gap by proposing models and tools that help in the assessment of LED-to-camera communication network mechanisms. Our contributions are threefold. First, we propose an analytical model for LED-to-camera communication systems, based on the theory of Markov-modulated Bernoulli processes (MMBP) [14], and show that it is significantly more accurate than a classical Gilbert-Elliott model [15]. Second, to facilitate performance evaluation tasks, we design, implement and validate *CamComSim*, the first LED-to-camera communication simulator. Finally, the model and the simulator allow us to benchmark several network redundancy mechanisms proposed in the literature. By checking against experimental results, we are able to confirm the correctness of our models in this context.

The remainder of this study is structured as follows. We discuss related works in Sec. 2 and we describe the testbed used for experiments in Sec. 3. We detail our analytical modeling of the LED-to-camera propagation channel in Sec. 4, complemented by a model of the receiver in Sec. 5. The *CamComSim* simulator is presented, validated and tested in Sec. 6, before concluding remarks in Sec. 7.

2. Related Works

The use of visible light for communications actually predates the one of RF. In fact, the first phone call in history was actually made using a photophone [16] and optical communications are a major component of telecommunication systems nowadays, through their use in optical fibers [17].

The use of wireless VLC gained a lot of interest in the last two decades, with the development of wireless local area networks (WLAN), as a competitor to the WiFi technology [18]. The term *LiFi* was coined in this sense [19] and numerous works discussed the advantages and disadvantages of this approach [5]. However, the use of VLC in WLANs requires dedicated receivers, usually incorporating sensitive photodiodes [1, 2]. In this sense, CoLight [20] interfaces a photodiode with a smartphone, using the audio jack and obtaining a throughput of 80 kbits/s.

In 2012, Danakis *et al.* [3] demonstrated for the first time that the rolling shutter effect of smartphone cameras can be used to receive information transmitted through visible light, creating the field of optical camera communications [21]. Using the smartphone camera as a receiver can also be used for other services. For example, *LiTell* [10] is a localization scheme that employs unmodified fluorescent lights (FLs) as location landmarks and relies on the observation that each FL has an inherent characteristic frequency which can serve as a discriminant feature, providing sub-meter accuracy. Using a similar idea, *iLAMP* [11] further introduces a sensor-assisted photogrammetry technique to estimate the 3D location with a small 90-percentile error of 3.5 cm. Another example is that of privacy, where *LiShield* [13] leverages the rolling shutter artifact to protect a physical scene from photographing, by illuminating it with smart LEDs transmitting modulated light.

However, the main usage of the rolling shutter effect remains LED-to-camera communication, which raises a lot of interest because it enables short-range communication with no extra financial cost between any LED-equipped machine or instrument and any regular smartphone. Both LOS and NLOS communication has been demonstrated, and a real competition in terms of throughput began. Lee *et al.* [6] extensively studied the rolling shutter effect and found an unpredictable and varying jitter between two consecutive frames. Their solution, *RollingLight*, uses a simple coding approach and provides a throughput of 20 bytes/s. Other solutions, such as *CeilingCast* [7] leverage rateless codes to face the unavoidable packet erasure caused both by the camera and the distance, reaching a throughput of 1.35 kbits/s. *MARTIAN* [8] exploits NLOS reflections and an original modulation technique to achieve a throughput of 1.6 kbits/s. In our own previous work [22], we relied on random linear coding to reach a throughput of 2.2 kbits/s. Relying on colored LEDs, *ColorBars* [23] demonstrates a throughput between 2.9 kbits/s and 7.7 kbits/s, depending on the receiving smartphone.

However, despite this significant interest in the subject, practically every

study on this topic uses an experimental approach. While experiments are essential in evaluating new protocols and services, running an experimental campaign every time one wishes to evaluate a new idea can become cumbersome. It is also very difficult to reproduce other solutions and fairly compare different techniques. Several VLC testbeds have been proposed in this sense [24, 25], but non-experimental performance evaluation tools in this context are poorly studied. A rare example is a channel model recently proposed by Zhang *et al.* [26], where the light attenuation between the LED and the camera is modeled depending on different physical parameters.

Designing analytical models and implementing them in simulation tools is standard practice in the wireless networking field in order to accelerate the evaluation of new protocols and mechanisms. Indeed, network simulation has been largely studied in the case of wireless communication [27]. Nonetheless, VLC performance evaluation remains poorly investigated and VLC simulation tools are still missing. The main efforts on simulating VLC systems have focused on indoor channel simulation [28], or on the 802.15.7 PHY [29, 30] and MAC [31] layers. These approaches rely on classical network simulation frameworks, such as those used for wireless and ad hoc networks, e.g. ns-2 [31], ns-3 [29], OMNET++ [30] or MATLAB [28]. All these works consider LED-to-Photodiode communication, hence OCC is completely unexplored. Our work is the first effort in LED-to-Camera simulation reported in the literature, making *CamComSim* the first implementation of a LED-to-Camera VLC simulator.

3. Testbed description

The analytical and simulation models proposed in this paper are compared and validated using an extensive experimental campaign. These experiments were conducted using a testbed detailed in our previous work [32]. We briefly describe its main components below.

We designed a 40x20 mm printed circuit board (PCB) shown in Fig. 2 that supports a red 5 mm LED, a micro-controller unit (MCU), a surface mount device (SMD) RGB LED whose color can be easily changed, and a temperature sensor, to broadcast the ambient temperature through the light. To evaluate the system, we also use a Nucleo development board to interface the MCU with more convenience, and so, change without hardware modification the LED type, the General Purpose Input/Output (GPIO) mapping or the firmware implementation. The development board is depicted in Fig. 3.

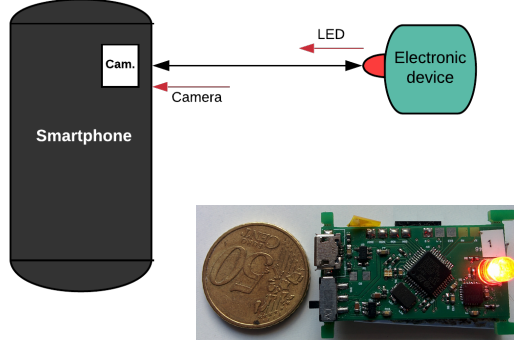


Figure 2: On the left, our LED-to-Camera VLC system. On the right, our smart VLC device prototype.

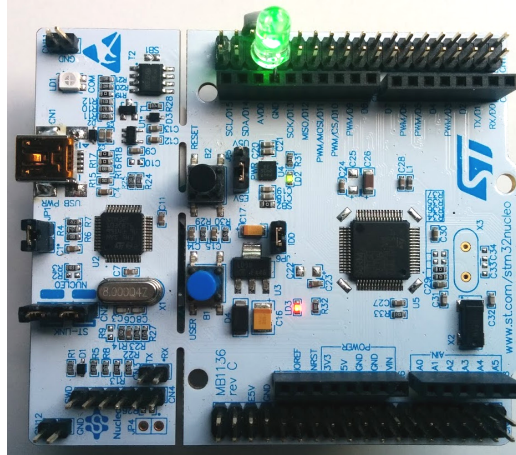


Figure 3: The Nucleo STM32L0 development board.

142 We choose the STM32L051 low cost and low power MCU from ST Micro-
 143 electronics, an MCU similar to those already integrated in most household
 144 appliances. The core is a Cortex M0+, running up to 32 MHz, with 32 Ko
 145 Flash and 8 Ko RAM. To get a better clock accuracy and avoid clock bias
 146 due to the temperature, we use an 8 MHz high speed external crystal oscilla-
 147 tor as the clock source and make the core run at this speed. As proposed in
 148 previous studies [3, 6], the LED signal is modulated using the on-off keying
 149 modulation scheme. We consider a clock-rate varying from 2 to 10 KHz,
 150 which is a suitable bandwidth for OCC when using the rolling shutter effect

151 [9]. To ensure a balanced duty cycle signal and avoid any flickering effect,
 152 we use the Manchester coding proposed in [3, 9].

153 On the receiver side, we use a LG Nexus 5 smartphone running Android
 154 Marshmallow version number 6.0.1. It has a Qualcomm Snapdragon 800
 155 quad-core CPU 2,26 GHz CPU and 2 Go RAM. Its 8 megapixels 1080p
 156 1/3.2" CMOS sensor with 1.4 μm pixel size can capture up to 30 frames per
 157 second and supports advanced imaging application provided by the Camera2
 158 API. We have developed an Android application that sets up the camera
 159 parameters to observe the rolling shutter effect produced by the modulated
 160 LED. For that, based on [9], we set a very short exposure time and an
 161 increased sensor sensitivity, respectively to 100 μs and ISO 10000. As soon
 162 as a new frame is available, the application creates and starts a new thread
 163 to process and decode the picture on the background.

164 4. Modeling LED-to-camera communication

165 In the following, we describe an OCC communication channel model.
 166 Based on the theory of Markov-modulated Bernoulli processes (MMBP) [14],
 167 discussed in Sec. 4.1, this model can be applied to all LED-to-camera com-
 168 munication systems. The proposed model is not only generic, but also very
 169 accurate, as demonstrated by its validation with extensive experimental re-
 170 sults in Sec. 4.4. As an example, we use the analytical model to compare two
 171 simple redundancy mechanisms, required to cope with the inherent losses of
 172 LED-to-camera communication and described in Sec. 4.2 and Sec. 4.3.

173 4.1. Model design

174 In a LED-to-camera system, data is received as a series of dark and
 175 illuminated stripes appearing in a region of a picture frame captured by the
 176 camera. We denote this part of the picture as a region of interest (ROI). In
 177 the following, we note by f_i the i -th frame captured by the camera and by δ_f
 178 the time between the beginning of two consecutive frames. Obviously, even
 179 at the highest frame rate allowed by the camera, data is not continuously
 180 received, as a minimum time δ_g exists between two frames. This is denoted
 181 as the inter-frame gap (IFG). Moreover, as depicted in Fig. 4, the distance
 182 between the LED and the camera also has an impact: when the camera is
 183 farther away, the LED transmission is captured for a shorter time, resulting
 184 in a smaller ROI.

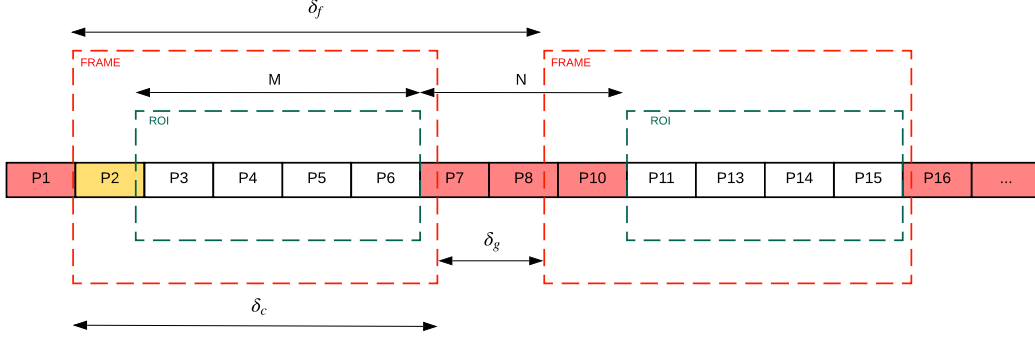


Figure 4: Frame capture time and inter-frame interval, and their relation with the MMBP parameters.

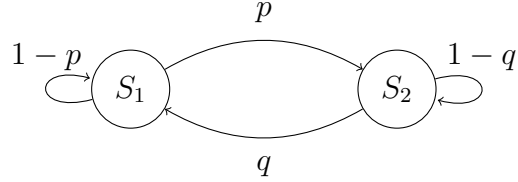


Figure 5: The *Gilbert-Elliott* model of the LED-to-camera channel.

185 4.1.1. *Gilbert-Elliott model:*

186 A first idea to model LED-to-camera communication would be the Gilbert-
 187 Elliot model, which is widely used to model bursty losses [15]. This unique
 188 type of channel can intuitively be modeled by a two states Markov chain,
 189 as depicted in Fig. 5. In state S_1 , the system is capturing a frame. The
 190 camera is receiving packets, and the reception probability is $1 - p_e$, where
 191 p_e is the packet decoding error probability. In state S_2 , the camera is not
 192 capturing any pictures, therefore we consider the reception probability is 0.
 193 The transition probability between S_1 and S_2 (respectively S_2 and S_1) is
 194 denoted p (respectively q). The model assumes that p, q, p_e are independent
 195 and constant.

196 In this case, the probability of being in state S_1 under the steady-state
 197 regime can be easily computed as $p_{s_1} = \frac{p}{p+q}$. The probability of being in
 198 state S_2 is so $p_{s_2} = \frac{q}{p+q}$.

199 The values of p and q are function of the duration of the IFG, δ_g , the
 200 frame duration, δ_f , and the camera capture time δ_c , depicted in Fig. 4, and
 201 linked as the following:

$$p = 1 - q = \frac{\delta_f - \delta_g}{\delta_f} = \frac{\delta_c}{\delta_f} \quad (1)$$

202 *4.1.2. MMBP model:*

203 If the model introduced just before is straightforward and widely used,
 204 it lacks realism in our case, where the transitions between ON and OFF
 205 states are almost deterministic. Practically, in our system, the transition
 206 probability from a state to another depends on the residence time in this
 207 state.

208 To improve this approach, we model the LED-to-camera channel using
 209 a Markov-modulated Bernoulli process (MMBP), represented in Fig. 6. In
 210 this figure, we depict a Markov chain with a total number of $M + N$ states.
 211 Each of these states represents a reception time slot, i.e. the time duration
 212 needed in order to receive one physical layer message (denoted as PHY-SDU
 213 in the following). The transition between two states representing successive
 214 time slots is automatic, i.e. it happens with a probability of 1.

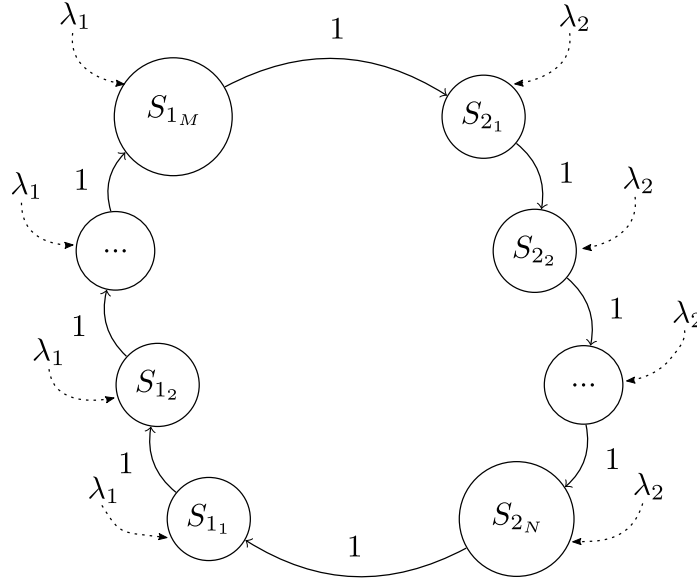


Figure 6: The MMBP model of the LED-to-camera channel.

215 Practically, the $M + N$ states in Fig. 6 represent a δ_f time interval,
 216 and they are divided in two groups: M states corresponding to the camera

capture time δ_c (S_{ON} states), and N states corresponding to the inter-frame time δ_g (S_{OFF} states). A Bernoulli arrival process is associated with each of these $M + N$ states, representing the reception of a packet.

In S_{ON} states, the camera is receiving packets, and the arrival rate is $\lambda_1 = (1 - p_e)$, where p_e is the packet decoding error probability. In S_{OFF} states, the camera is not capturing any pictures, therefore we consider the arrival rate $\lambda_2 = 0$.

We denote as s a state in the Markov chain and we define state $s + j$ as the state reached after j transitions, starting from state s . The probability of being in state s under the steady-state regime can be easily computed as $\pi_s = \frac{1}{N+M}$. At the same time, the probability of noticing no arrivals (i.e. no packet reception) in state s is $p_0(s)$. This can be written as:

$$p_0(s) = \begin{cases} 1, & \text{if } s \in N \\ p_e, & \text{if } s \in M \end{cases} \quad (2)$$

As it can be seen from both models, the relatively high packet loss probability (compared with RF technologies) is an intrinsic property of the LED-to-camera communication channel. To overcome this problem, redundancy mechanisms are needed.

In the following, using classical redundancy mechanisms as an example, we show that the classical Gilbert-Elliot model is inaccurate, which highlights the need to rely on the MMBP theory. Then, we use the MMBP channel model to compare two simple, but widely used redundancy solutions: repeating a packet or repeating a sequence of packets.

4.2. Repeat Packet

The first strategy to cope with the inherent losses in the LED-to-camera communication system, used for example by Ferrandiz-Lahuerta *et al.* [9], is to send each packet twice in a row, to increase the probability that at least one of the transmissions will be fully captured by the smartphone camera. We generalize this approach in the Repeat Packet (RP) strategy, where each packet is repeated r times, one after the other. In this case, the r value needs to be chosen in order to attain a desired reception probability, its optimal value depending on the inter-frame time and on the packet size.

In the following, we study the probability of receiving a packet at least once when considering the RP strategy, for the two models introduced above.

244 *4.2.1. Gilbert-Elliott Model:*

If we consider the Gilbert-Elliott model, the probability of receiving a packet at least once can be written as $p_s^{RP} = 1 - p_0^{RP}$, where p_0^{RP} represents the probability of failing to receive a packet r times in a row, written as:

$$p_0^{RP} = (P_{S_1} \cdot p_e + P_{S_2})^r = \left(\frac{p \cdot p_e + q}{q + p} \right)^r \quad (3)$$

245 *4.2.2. MMBP Model:*

If we consider the MMBP model, the probability of failing to receive a packet r times in a row p_0^{RP} can be written as:

$$p_0^{RP} = \frac{\sum_s (p_0(s) \cdot p_0(s+1) \cdot p_0(s+2) \cdot \dots \cdot p_0(s+r-1))}{N+M} \quad (4)$$

The value of p_0^{RP} will depend on m_s , defined as the number of S_{ON} states during the r retransmissions, when the first packet transmission is in state s :

$$p_0^{RP} = \frac{1}{N+M} \cdot \sum_{s=1}^{M+N} (p_e)^{m_s} \quad (5)$$

246 Depending on the values of r , M and N , several cases can be distin-
247 guished. We present results for the two most current cases:

Case 1: $r < M, N$. This means that the number of retransmissions does not always cover the δ_g period (which counts N states). In this case, p_0^{RP} can be written as follows:

$$\begin{aligned} p_0^{RP} &= P[s \in N \wedge (s+r) \in N] + P[s \in M \wedge (s+r) \in M] \\ &\quad + P[s \in M \wedge (s+r) \in N] + P[s \in N \wedge (s+r) \in M] \\ &= \frac{N-r+1}{M+N} + \frac{M-r+1}{M+N} \cdot p_e^r + \frac{1}{M+N} \sum_{i=1}^{r-1} p_e^i + \frac{1}{M+N} \sum_{i=1}^{r-1} p_e^{r-i} \\ &= \frac{N-r+1}{M+N} + \frac{M-r+1}{M+N} \cdot p_e^r + \frac{2}{M+N} \cdot \frac{p_e - p_e^r}{1 - p_e} \end{aligned} \quad (6)$$

Case 2: $N < r < M$. . This is the most common case, where the number of repetitions is chosen to cover the entire inter-frame period. However, a

reception is still not certain in this case, because of the decoding error p_e . In this case:

$$\begin{aligned}
p_0^{RP} &= P[s \in M \wedge (s+r) \in M] + P[s \in M \wedge (s+r) \in N] \\
&\quad + P[s \in N \wedge (s+r) \in M] \\
&= \frac{M-r+1}{M+N} \cdot p_e^r + \frac{1}{M+N} \sum_{i=1}^{r-1} p_e^i + \frac{1}{M+N} \sum_{i=1}^N p_e^{r-i} \\
&= \frac{M-r+1}{M+N} \cdot p_e^r + \frac{1}{M+N} \cdot \frac{p_e - p_e^r}{1 - p_e} + \frac{1}{M+N} \sum_{i=1}^N p_e^{r-i} \\
&= \frac{M-r+1}{M+N} \cdot p_e^r + \frac{1}{M+N} \cdot \frac{p_e - p_e^r}{1 - p_e} + \frac{1}{M+N} \cdot \frac{p_e^{r-N}(p_e - p_e^N)}{1 - p_e}
\end{aligned} \tag{7}$$

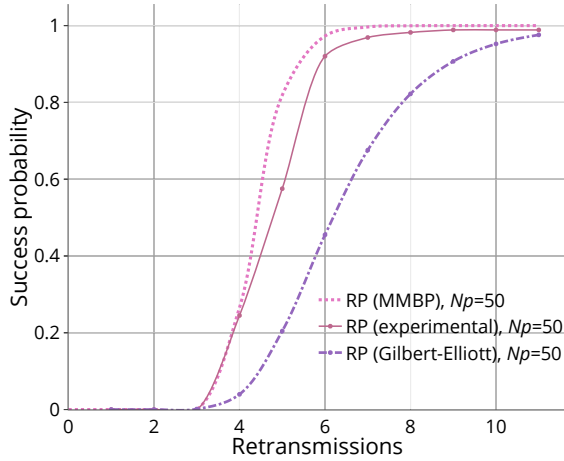


Figure 7: Probability to successfully receive N_p packets for the RP strategy. Dotted-lines show analytical results for the Gilbert-Elliott and MMBP models, while plain lines represent experimental results.

248

249 We compare the analytical results given by the two aforesaid models to
250 experimentation results, obtained using our testbed, in Fig. 7. This figure
251 shows the success probability of receiving a message of $N_p = 50$ packets of
252 data, as a function of the number of retransmissions r .

253 The very different results between the Gilbert-Elliott model and the ex-
254 perimentation confirms that the stochastic transition assumptions of this

255 model are quite far from reality. On the other hand, MMBP approximates
 256 quite well the experimental behavior, highlighting the need for this more
 257 complex, but finer grained model.

258 4.3. Repeat Sequence

259 A different approach to improve reliability is the Repeat Sequence (RS)
 260 strategy, consisting in the transmission of a sequence of N_p packets, repeated
 261 r times. In contrast with the previous mechanism, RS does not try to cover
 262 the inter frame time at the packet level, and it does not ensure that a packet
 263 is received before sending the next one. Instead, the reliability and presumed
 264 efficiency is based on the fact that the probability of losing the same packets
 265 over different transmitted sequences is low.

In the case of an RS strategy with a sequence of N_p packets retransmitted
 r times, the probability of receiving a packet at least once can be written as
 $p_s^{RS} = 1 - p_0^{RS}$. Using the MMBP model, the probability of failing to receive
 a packet r times in a row, p_0^{RS} , can be written as:

$$p_0^{RS} = \frac{1}{M+N} \cdot \sum_{s=1}^{M+N} \prod_{j=0}^{r-1} p_0(s + rN_p) \quad (8)$$

266 4.4. Evaluation results

267 We use our MMBP analytical model to study the RP and RS strategies
 268 by focusing on the probability of delivering the entire quantity of information
 269 in a given number of transmissions. We provide both analytical and
 270 experimental results, allowing us to validate the proposed MMBP model.

271 Fig. 8 shows, for the two mechanisms, the probability of integrally re-
 272 ceiving N_p packets of data as a function of the number of retransmissions
 273 r . In this figure, we set $M = 5$ and $N = 2$; these values are in line with
 274 the packet length, the transmitter frequency and the camera capture interval
 275 experimentally observed for a distance of 5 cm between LED and camera.
 276 The results show quite a nice fit between the analytical and experimentation
 277 results, despite the assumptions required by our MMBP model.

278 To better understand the performance of the two retransmission strate-
 279 gies, we compare them in Fig. 9. This figure shows that, for the RS strategy,
 280 3 retransmissions are needed to achieve a reception probability higher than
 281 0.9, while this value raises to 6 for the RP strategy. On the right side of the
 282 figure, we show that the performance of the two strategies depends on the
 283 ratio between the number of S_{ON} and S_{OFF} states, $M : N$. When this ratio

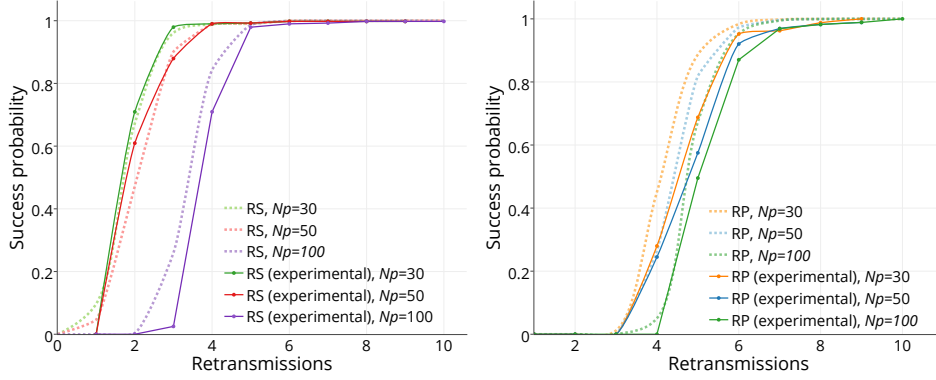


Figure 8: Probability to successfully receive N_p packets for the RS (left) and RP (right) strategies as a function of the number of retransmissions r . Dotted-lines show analytical results while plain lines represent experimental results.

changes from 5 : 2 to 2 : 5, which practically corresponds to increasing the distance between the LED and the camera, RP gives better results than RS. Indeed, for the RS method, the success probability sharply decrease when $M < 3$ and stays below 0.6 even for 10 retransmission.

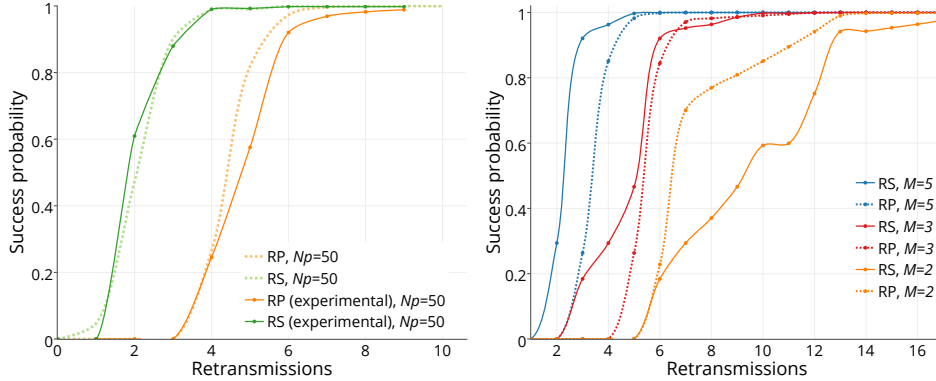


Figure 9: Comparison between RS and RP. On the left, analytical and experimental results for $M = 5$ and $N = 2$. On the right, analytical results when $M + N = 7$, but the $M : N$ ratio changes. In both cases, $N_p = 50$.

Practically, this means that RP is more suitable when the distance between the LED and the camera is higher, while RS is better for short communication distances. This phenomenon was previously unknown in the research community, but it is straightforward to study with our analytical model

292 5. ROI model

293 An important phenomenon in LED-to-camera communications comes as a
 294 direct consequence of the distance between the LED and the camera. Indeed,
 295 as this distance increases, the size of the ROI in the picture reduces and, as a
 296 consequence, cuts down the number of messages that the camera can receive
 297 per frame, i.e. the M states in Fig. 4. To include this performance factor into
 298 our model, we propose an analytical function that gives the ratio between
 299 the ROI and the picture size. In the model discussed in Sec. 4.1, this is the
 300 ratio of M states in the $M + N$ states.

301 We apply photogrammetry rules to give the ROI ratio as a function of the
 302 distance d , the LED size l , the camera CMOS sensor size ss , the image size
 303 on the sensor i and the camera focal distance fc . According to the optical
 304 system depicted in Fig 10, basic lens optic rules give the following Eq. 9:

$$\frac{i}{fc} = \frac{l}{d} \iff i = \frac{l \cdot fc}{d} \quad (9)$$

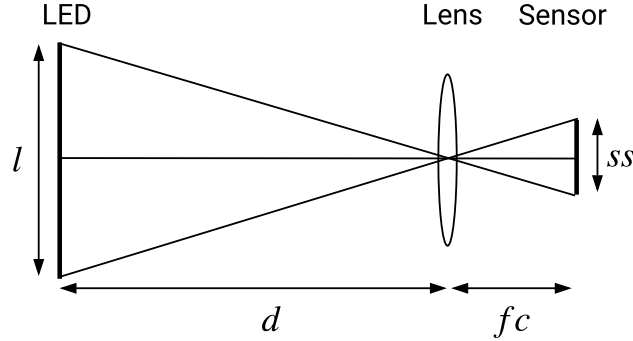


Figure 10: Formation of an image on a sensor by a converging lens.

To obtain the ROI as the ratio of the total number of pixels in the picture, we need as input the CMOS sensor size ss . We apply the $\min()$ function to normalise the $ROI \in [0, 1]$ even if the image size on the sensor, i , is larger than the sensor size, ss . The result is given in the following Eq. 10:

$$ROI = \min \left(1, \frac{l \cdot fc}{d \cdot ss} \right) \quad (10)$$

305 To validate the results given by Eq. 10, we measure the ROI experimen-
 306 tally, using the testbed described in Sec. 3, for distances from 0 to 40 cm.

307 Fig. 11 plots in orange the ROI ratio we observed during our experiments
 308 and in green the analytical results computed with a Nexus 5 sensor with the
 309 following characteristics: $fc = 35$, $ss = 5.7$ and $l = 10$. This shows that
 310 the analytical curve approximates quite well the experimental ROI ratio.
 311 However, we notice that the experimental results are better for a distance
 312 between 10 and 30 cm, and they become worse than the model at 35 cm. In
 313 fact, the light radiance on the camera lens, that our model does not take into
 314 account, artificially increases the LED size on the picture when the camera
 315 is close to the LED. The difference at larger distance is a consequence of the
 316 ambient light which was measured at 650 lux during the experiments, also
 317 neglected in Eq. (10).

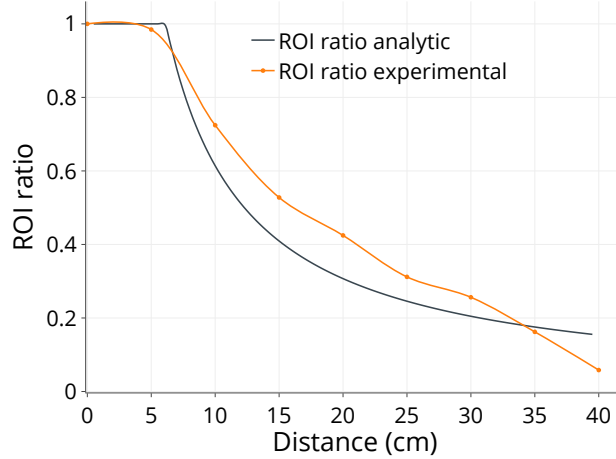


Figure 11: ROI as a function of distance. The orange line shows experimental results, while the green line represents analytical results given by Eq. 10.

318 The ROI model described in this section and the MMBP reception model
 319 validated in the previous section are the basis of the simulator implementa-
 320 tion discussed in the following.

321 6. The CamComSim Simulator

322 As discussed in Sec. 2, the simulation of LED-to-camera communication
 323 remains completely unexplored in the field. Our work is the first such effort
 324 reported in the literature, making *CamComSim* the first implementation of
 325 a LED-to-camera VLC simulator.

6.1. Motivation

The channel and receiver models described in the previous sections allow us to obtain very accurate results, closely matching the experimental values. However, the two models are deterministic by nature, and they can lead to strange artifacts in some specific cases.

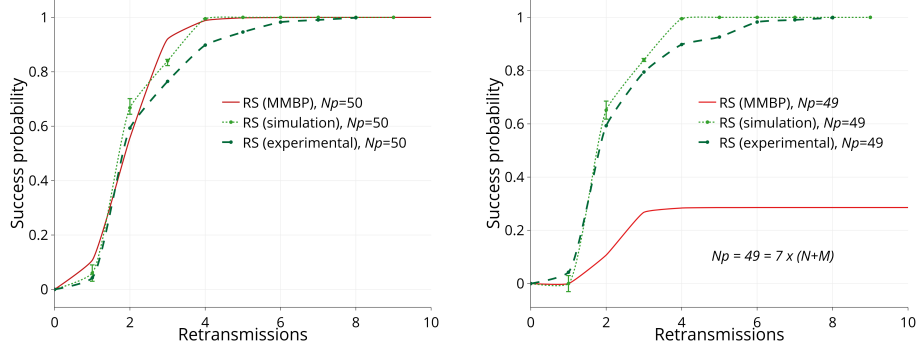


Figure 12: Analytical, simulation and experimental results obtained for $N_p = 50$ (left) and $N_p = 49$ (right).

One such example is provided in Fig. 12, where analytical and experimental results are compared for two different numbers of transmitted packets, $N_p = 49$ and $N_p = 50$. While the analytical model is very accurate for $N_p = 50$, things are very different for $N_p = 49$, where the MMBP model indicates that less than 30% of the packets should be received, regardless the number of retransmissions. This is an artifact of the MMBP model, which considers the time between two consecutive pictures taken by the camera to be perfectly constant. Because of this, when the number of packets N_p is a multiple of the number of states $M + N$ in the MMBP model, a repetitive pattern appears and all the retransmissions of some packets should be lost, resulting in a maximum packet reception probability of $N/(M + N)$. However, the experimental results indicate that this property is not actually present in real systems, where the inter-frame gap sufficiently variable to remove this phenomenon. Accounting for this analytically would be possible, but it would require an even more complex modeling approach.

While built on the models described above, a simulation approach can address this problem in a much simpler way. Furthermore, a simulator can easily account not only for inter-frame variability, but also for more complex events, such as user motion. This is demonstrated in Fig. 12, where

the results obtained with *CamComSim* closely match those obtained on the testbed.

6.2. Simulator Implementation

6.2.1. Software architecture

CamComSim is an event-driven LED-to-camera simulator developed in Java, which makes it easy to maintain and distribute code, and it provides built-in multi-platform compatibility for systems with a Java Virtual Machine. Fig. 13 shows the *CamComSim* software architecture that consists of a simulator kernel class and four core packages. For interested readers, *CamComSim* is already available as an open-source software under Apache license at <http://vlc.project.citi-lab.fr/camcomsim>.

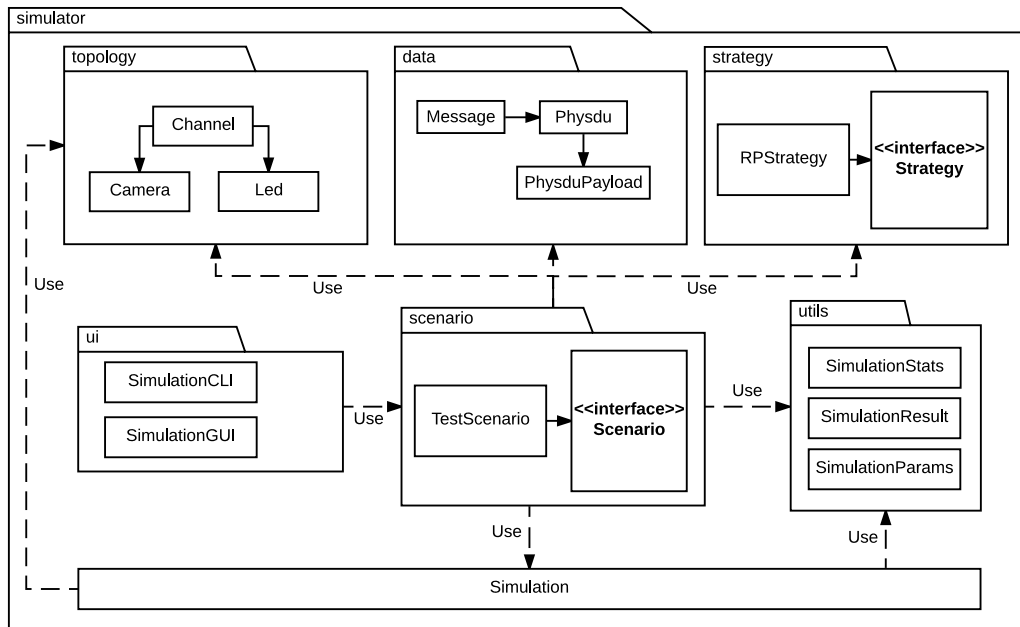


Figure 13: The *CamComSim* software architecture and packages dependency graph.

The **topology** package groups classes that describe the system components: **Led**, **Camera** and **Channel**. The classes in the **data** package implement the data encapsulation. For this, a **Message** is a set of PHY-SDU that encapsulates a **PhysduPayload**. A **Packet** is a **PhysduPayload** child class, with a sequence number as header and a payload that contains

366 data. Before each simulation, a `Message` is created according to the user
 367 settings. The resulting set of `Physdu` is initialized with a `Packet` filled
 368 with arbitrary data in the payload (real data could be used if available) and
 369 a unique sequence number in the header.

370 The broadcast strategy abstraction is given in the `strategy` package.
 371 Here, the `Strategy` interface lets the users implement their transmission
 372 strategy. This package also contains the straightforward implementation of
 373 the RP strategy, `ReapeatPhysduStrategy`, which consists in repeating each
 374 PHY-SDU r times, one after the other, as described in Sec. 4.2. When the
 375 last PHY-SDU of the message is reached, the process is repeated from the
 376 beginning.

377 Finally, the `scenario` package proposes an interface to build a simula-
 378 tion, wiring together the `Message`, the `Channel`, the `Led`, the `Camera` and
 379 the `Strategy` with the `Simulator` kernel. Besides, the package `utility`
 380 provides helper classes used to compute the simulation results statistics, for-
 381 mat and save the results as a JSON file and load or save the simulation
 382 parameters. The `ui` package contains a command-line interface (CLI) used
 383 to run a simulation scenario.

384 6.2.2. Simulator parameters

385 Our simulator exposes a set of finely grained parameters to describe the
 386 LED-to-camera communication system behavior. Table 1 shows the param-
 387 eters we use in this work and expose through the *CamComSim* CLI. The
 388 performance of the LED-to-camera communication is significantly affected
 389 by the distance d , the IFG noted δ_g in Fig. 6, and the LED size l . The values
 390 of these parameters should therefore be carefully chosen, according to the
 391 available hardware and envisioned scenario.

392 Further parameters, introduced in Sec. 5, that refer to the CMOS sen-
 393 sor characteristics, are optional but can be considered to refine the channel
 394 model, as they impact the ROI. However, smartphone manufacturers rarely
 395 provide this information, e.g. regarding the sensor size ss and the focal
 396 distance fc .

397 The PHY-SDU error rate (PER) p_e is the consequence of the errors oc-
 398 ccurring in a M state when a PHY-SDU is well included in a picture but
 399 is wrongly decoded by the smartphone. These errors are bits substitutions
 400 induced by interference, low SNR, and artifacts on the picture. This value
 401 varies from a smartphone to another, but we did not observe major differ-

Param.	Description	Default Value
d	Distance between camera and LED (cm)	5
l	LED size (mm)	4
d_g	Camera inter-frame gap ratio	0.1
p_e	Decoder PHY-SDU Error Rate	0.001
f	Modulation frequency (Hz)	8000
P	PhysduPayload Header length (bit)	8
H	PhysduPayload Payload length (bit)	16
r	PHY-SDU repeat number	1
G	Message size (bytes)	50

Table 1: Simulator parameters and default values.

ences in our tests.

The **PhysduPayload** payload size P and the **PhysduPayload** header size H configure the data encapsulation mechanism. Given these two settings, the PHY-SDU size is computed as $P + H + SYNC + PB$, where $SYNC$ is the PHY-SDU delimiter symbol (4 bits in our tests), and PB is the number of parity bits (2 bits in our settings). This PHY-SDU size, along with d , l , δ_f and the modulation frequency f , determines the number of the M time slots in Fig. 6.

A transmission strategy among the **Strategy** interface implementations needs to be chosen as well. For now, we have implemented and considered only the RP strategy, for which the parameters are the size of the message to broadcast, G , and the number of consecutive PHY-SDU emissions, r .

6.2.3. Kernel Implementation

The *CamComSim* kernel is implemented in the **Simulator** class. Its role is to produce PHY-SDU emission events (TX) and manage their result. The number of events, i.e. the number of PHY-SDU sent, is noted c and is updated at runtime. At each clock tick, c is incremented, a TX event is created and processed as follow: (1) the next PHY-SDU in the transmission strategy queue is associated with this event; (2) considering p_e , f , P , H , c , the channel response function gives the event result. The result is one among: *reception success*, *reception with errors* or *loss during IFG*; (3) this result is stored in a list to further determine if all the PHY-SDU forming a

424 message are received.

425 The simulator loops over (1), (2) and (3) until the stop condition is met,
426 i.e. c has reached the maximum number of PHY-SDU emissions or the com-
427 plete message is received. The simulation is repeated n_r times using the Java
428 class for multi-threading purpose `ThreadPoolExecutor`. Finally, the simu-
429 lations results and statistics are saved in a JSON file for further processing.

430 6.3. *CamComSim* validation

431 In this section, we present LED-to-camera simulation results given by
432 *CamComSim*. To assess the correctness of our simulator, we conduct a se-
433 ries of experiments with the testbed presented in Sec. 3. We set the emitter
434 symbol rate to 8 kHz and place it in standard indoor illumination conditions,
435 near a window and illuminated with neon lights. The illuminance has been
436 measured with a luxmeter at around 650 lux. We compare the testbed per-
437 formance with the results given by *CamComSim* for a set of key parameters:
438 the message size G , the number of consecutive PHY-SDU emissions r , the
439 distance d and the PHY-SDU payload length P .

440 6.3.1. PHY-SDU Retransmission

441 As discussed in Sec. 4.2, to face the IFG bits erasure and ensure that
442 all the packets are well received, a possibility is to transmit consecutively
443 each PHY-SDU r times in a row. The `ReapeatPhysduStrategy` class in
444 *CamComSim* implements this retransmission strategy.

445 Fig. 14 shows the goodput at 5 cm for different values of PHY-SDU
446 consecutive retransmissions r , with $G = 50$, $P = 19$, $H = 5$. When each
447 PHY-SDU has been transmitted r times, the message transmission restarts,
448 until the message is completely received. To avoid infinite loops, we stop
449 the simulation when 50000 PHY-SDU are sent, even if the message is not
450 received entirely. In such case, we consider the goodput is 0.

451 The results highlight that the simulation and testbed goodput follow the
452 same tendency when r varies. The best case is when $r = 1$, for which
453 the goodput is 1.6 kbit/s according to *CamComSim* and 1.7 kbit/s for the
454 testbed, an estimation error of only 6%. Based on these results, for all the
455 simulations that follow we use the RP strategy implementation with $r = 1$.

456 6.3.2. Message size

457 We now consider the impact of the message size G on the goodput at a 5
458 cm distance, with $r = 1$ and a 24 bits length PHY-SDU. Fig. 15 shows that

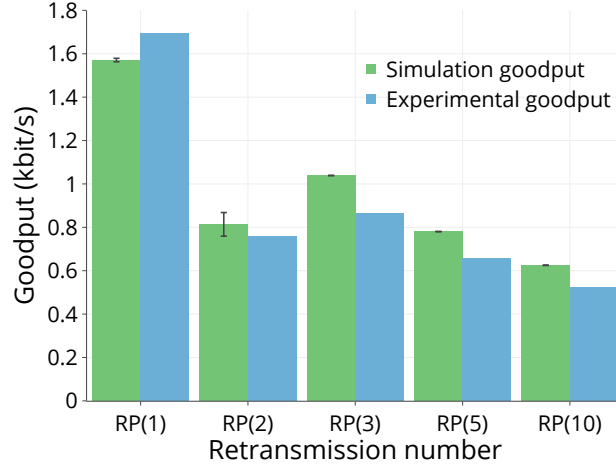


Figure 14: The experimental goodput (blue) compared with the simulation goodput (green) as a function of the number of consecutive PHY-SDU emissions. The bars on top are 95% confidence intervals.

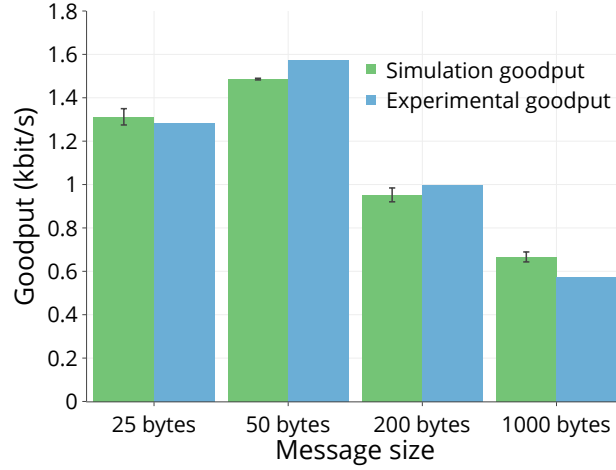


Figure 15: The experimental goodput (blue) compared with the simulation goodput (green) for different message size (G , bytes).

459 *CamComSim* results are very close to those of the testbed, confirming that
 460 the simulator well considers the impact of the message size. The goodput
 461 reduces when the message size increases, as the RP strategy leads to a large
 462 number of useless transmissions: the simulator gives 1.6 kbit/s of goodput
 463 for $G = 50$ bytes, while this falls to 670 bit/s when $G = 1000$ bytes. These

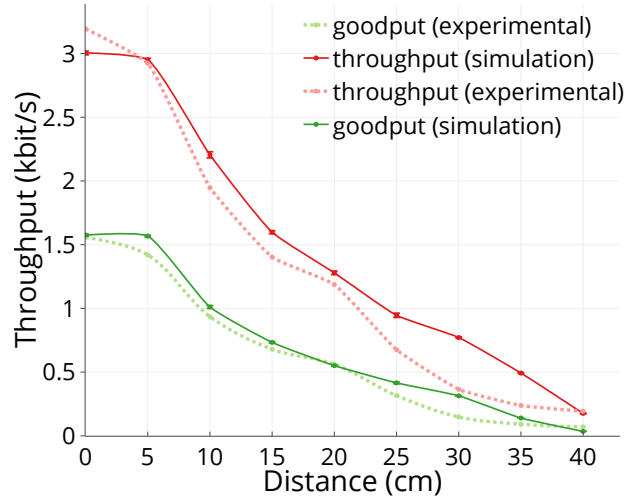


Figure 16: The throughput (red) as a function of the distance, compared with the goodput (green). Dotted-lines show experimental results while plain lines represent simulation results.

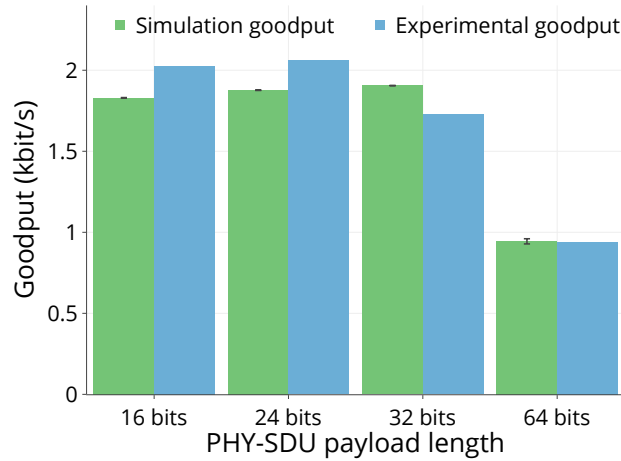


Figure 17: The experimental goodput (blue) compared with the simulation goodput (green) for different PHY-SDU payload size (bit).

464 results differ from the testbed in no more than 7%.

465 6.3.3. Distance

466 Fig. 16 shows the goodput and the throughput as a function of the dis-
 467 tance, when the LED broadcasts a 50 bytes message. The PHY-SDU payload

468 is set to 24 bits, with $P = 19$ and $H = 5$. The results show a good match
 469 between the simulation and real life results. At 10 cm, *CamComSim* gives
 470 2.2 kbit/s of throughput, against 1.94 kbit/s experimentally. The results are
 471 closer for the goodput: 0.94 and 1.0 kbit/s respectively for simulation and
 472 experimentation, that is only 6% of difference.

473 6.3.4. PHY-SDU length

474 Fig. 17 shows the impact of the PHY-SDU payload size on the goodput
 475 at 5 cm, with $G = 50$. The packets are built using the best value for P and
 476 H , that is to say with just enough bits in the header to label each packet
 477 with a unique sequence number. Both experimental and simulation results
 478 show that the optimal PHY-SDU payload size is 24 bits: *CamComSim* gives
 479 a goodput of 1.9 kbit/s, while the testbed reaches 2.1 kbit/s. These results
 480 outline that using large PHY-SDUs reduces the encapsulation overhead, but
 481 increases the probability that the IFG and a small ROI truncate the PHY-
 482 SDU. Fig. 17 brings out that *CamComSim* well considers this behavior: the
 483 goodput reaches 0.9 kbit/s for a PHY-SDU payload of 64 bits, very close to
 484 the experimental results.

485 Overall, this entire evaluation highlights that *CamComSim* gives results
 486 very close to the testbed for all the parameters we have studied. The differ-
 487 ence is around 10% and often less. For all the cases we consider, *CamComSim*
 488 respects the behavior of the LED-to-camera communication system imple-
 489 mented by the testbed.

490 6.4. Use case

491 In this section, we detail a case study for *CamComSim*, applied to a
 492 real life scenario. A common issue with cheap consumer electronics is the
 493 lack of diagnostics when a dysfunction happens. Manufacturers often blink
 494 the state LED with a pattern and color that match with an error code.
 495 Such a mechanism is easy to implement but leads to inaccurate diagnostics.
 496 For these cases, we propose to benefit from this LED to perform LED-to-
 497 camera communication and broadcast a log file that would include helpful
 498 information to diagnose a dysfunction. We consider a worst case file size of
 499 1 kbyte that is large enough for events history or debug traces.

500 Fig. 18 compares the goodput given by *CamComSim* with the goodput
 501 that our testbed achieved for the transfer of a 1 kbyte log file as a function
 502 of the number of PHY-SDU retransmissions r . Note that this is equivalent
 503 to $G = 1000$ bytes in Fig. 15. The transmission restarts until the message

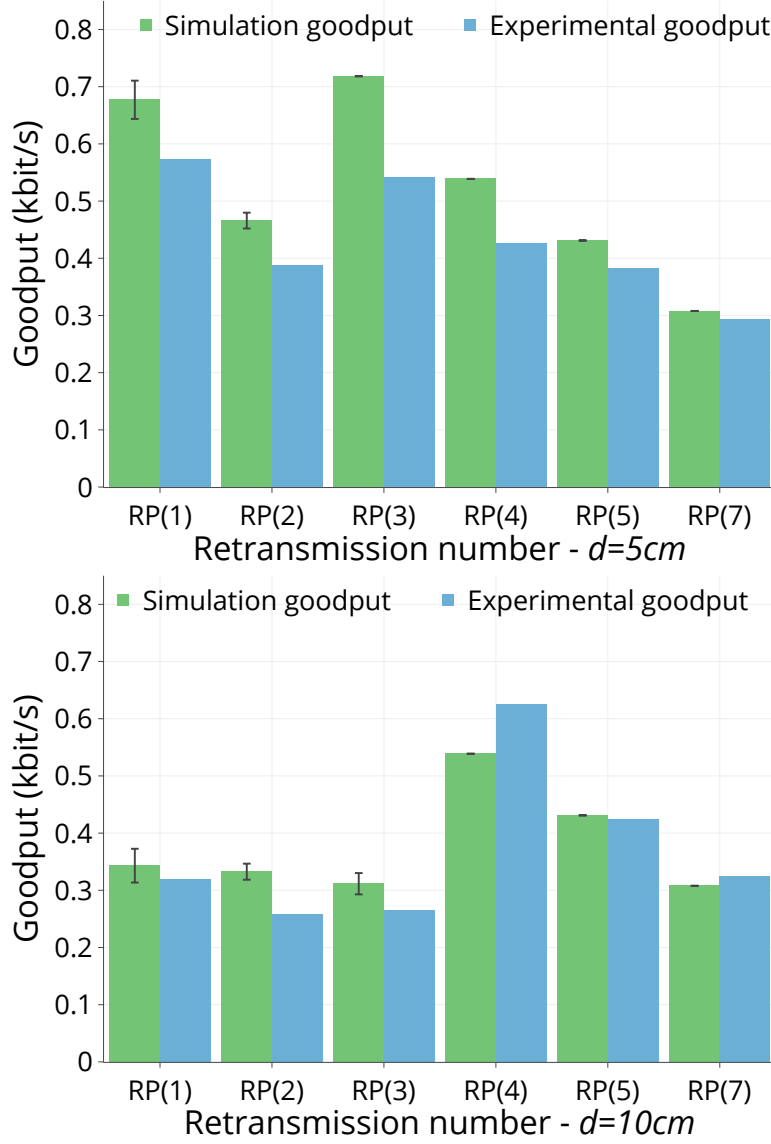


Figure 18: The experimental goodput (blue) compared with the simulation goodput (green) for the use case as a function of the number of consecutive PHY-SDU emission at 5 cm (left) and 10 cm (right).

504 is received. The left side plot shows the results when the LED and the
505 smartphone are 5 cm apart, while the distance is 10 cm on the right side
506 figure. At 5 cm, the simulation brings out that, to obtain the higher goodput,

the emitter should send each PHY-SDU one or three times consecutively, i.e. $r = 1$ or $r = 3$. The goodput is respectively 680 and 720 bit/s in these cases. This finding is similar to the testbed, where the goodput is 570 bit/s when $r = 1$ and 540 bit/s when $r = 3$.

Because the ROI decreases with the distance, the behavior is different when the smartphone is 10 cm far from the LED. In this situation, $r = 4$ stands out clearly to be the best choice both for the simulation and the experiments. The goodput then becomes 620 bit/s on the testbed and 540 bit/s with *CamComSim*.

Since the results are very close to the reality, using *CamComSim* highly reduces the search space for the experimental optimization of a system. As shown by these results, the best value for r can be decided using simulations only, removing the need for a lengthy experimental campaign.

7. Conclusion

In this paper, we introduced *CamComSim*, the first simulator for the design, the prototyping and the development of protocols and applications for LED-to-camera communication. Our event driven simulator is based on an MMBP channel model, and it relies on a standalone Java application that is easily extensible through a set of interfaces. We have validated *CamComSim* comparing simulation results with the performance reached by a real life testbed. Then, we illustrated with a practical use case the complete usage of *CamComSim* to tune a broadcast protocol that implements the transmission of a 1 kbyte log file. The results highlight that our simulator is very precise and can predict the performance of a LED-to-camera system with less than 10% of error in most cases. The availability of accurate performance evaluation tools offers a great ease of use and the opportunity to tune protocols without the burden of always realizing experiments on a testbed.

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